# TOSS: A Topology-based Scheduler for Storm Clusters

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# Agenda

- MOTIVATION
- SYSTEM DESIGN & IMPLEMENTATION
- EVALUATION AND EXPERIMENTAL RESULTS
- CONTACT INFORMATION

### MOTIVATION

- Inter-node traffic is likely to become a performance bottleneck.
- Scheduling executors represented in form of topologies play a vital role in optimizing performance of storm clusters.
- Reducing rescheduling overhead in the run time.

#### SYSTEM DESIGN TOSS Scheduler Partition Allocation Collect runtime Zookeeper data Worker Nodes Worker Node Supervisor Fig. 6: The System Design of the TOSS scheduler. **Executor Executor** Master Node **Executor** Nimbus Supervisor Executor **Executor** Fig. 2: A master node and worker nodes in Storm.

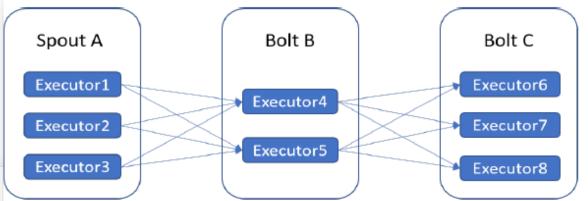


Fig. 4: Communications among multiple executors before executor allocation.

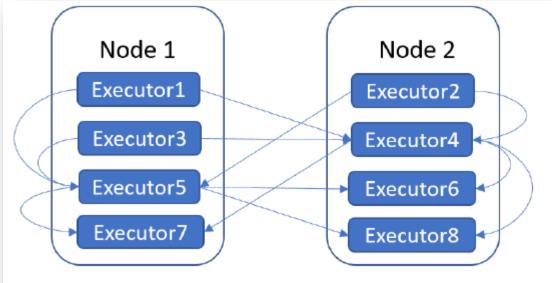


Fig. 5: Communications among multiple executors after executor allocation.

## SYSTEM DESIGN Cont. - Algorithms Design

phase

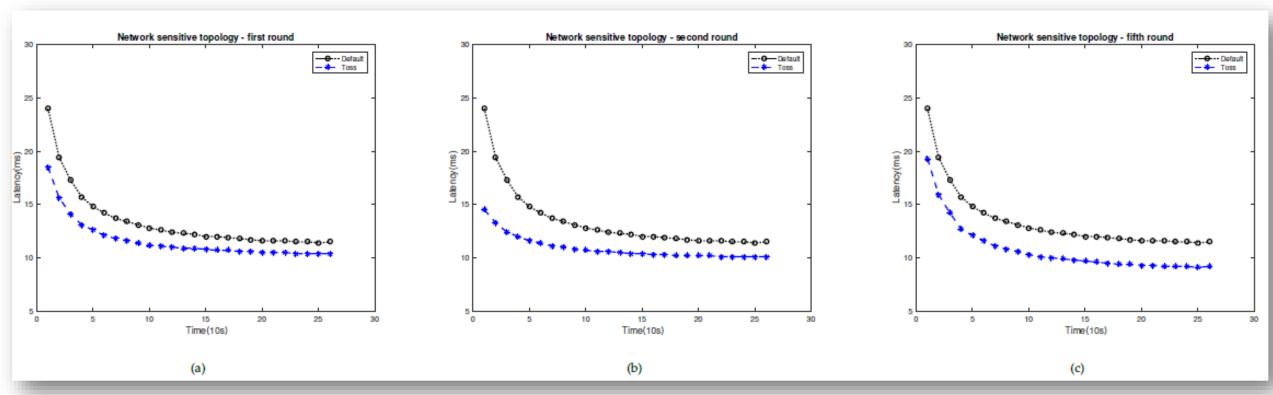
```
1 procedure TOSS;
   Input: Initial parameter set \alpha, topology set \Omega for
             allocation
   Output: Allocation assignment \Theta for each executors
2 Self-tuning parameter set \alpha;
E \leftarrow \text{total number of executor for allocation}
4 for topology \omega_i \in \Omega do
       for parameter \alpha_i \in \alpha do
           E_i \leftarrow E * \alpha_i;
           \theta_i \leftarrow \text{empty set};
           Runs DFS algorithm traverse, finds one chain.;
           C \leftarrow head of the chain:
           while \theta_i.size() < E_i do
               \theta_i add e_k \in E in C;
11
               C \leftarrow next component with executor for
                 allocation;
               if C reaches the end of chain then
                   C \leftarrow \text{back to head};
               end
           end
           Add \theta_i into \Theta;
           if C remains executor for allocation then
               Leave for next assignment computing
                 iteration;
           end
       end
22 end
```

```
Algorithm 1: TOSS scheduling algorithm: partition phase
```

```
procedure TOSS;
Output: Executor allocation on all nodes
Collects runtime workload data;
pq ← priorityQueue contains all current runtime workload for nodes N;
for assignment θ<sub>i</sub> ∈ Θ do
n<sub>i</sub> ← pq.pop();
assign θ<sub>i</sub> to node n<sub>i</sub>;
end
Algorithm 2: TOSS scheduling algorithm: allocation
```

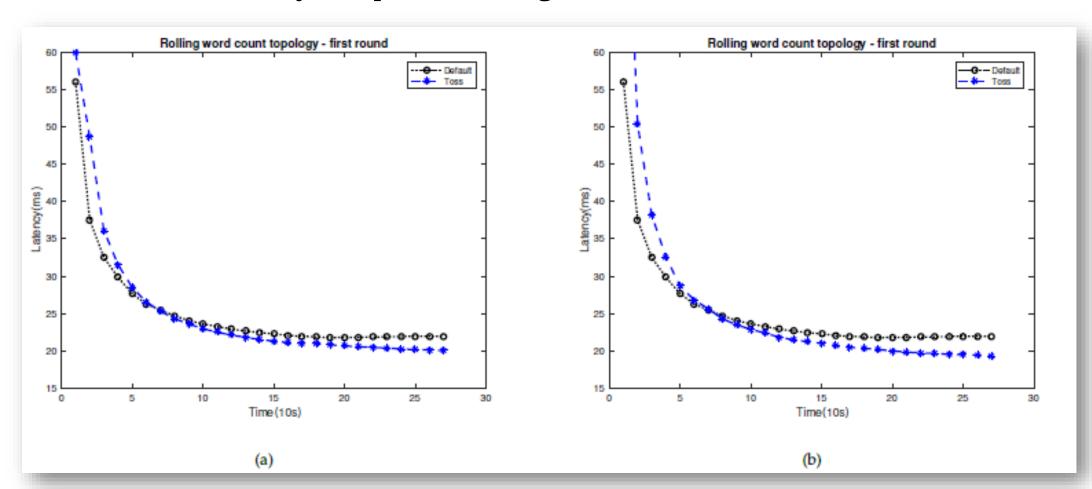
#### EVALUATION AND EXPERIMENTAL RESULTS

• The latency comparison between multiple submission of TOSS and default scheduler



#### EVALUATION AND EXPERIMENTAL RESULTS

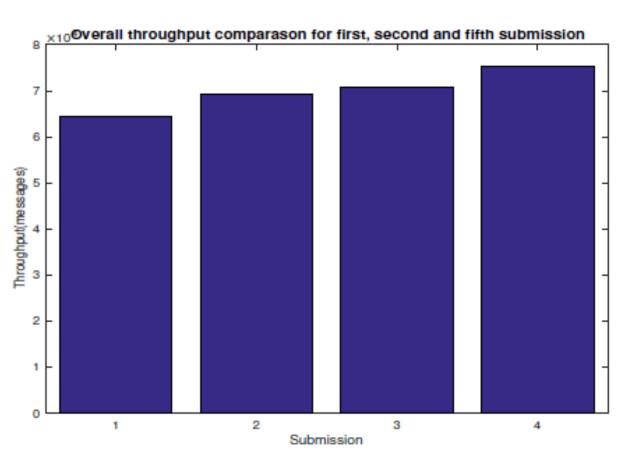
• The latency comparison using the Wordcount benchmark.

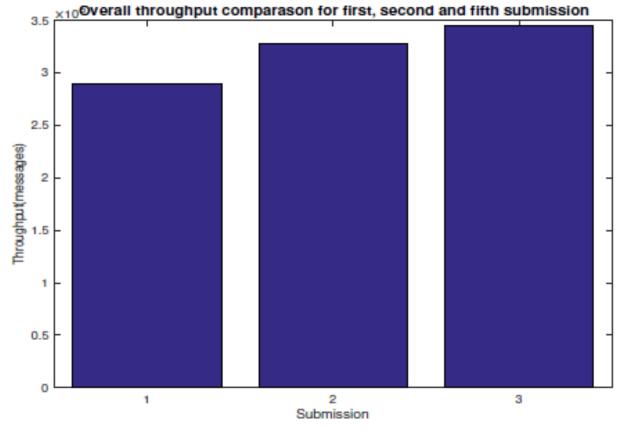


#### EVALUATION AND EXPERIMENTAL RESULTS

• The throughput comparison between default scheduler and multiple runs TOSS.

 The throughput comparison between TOSS and the exisiting scheduler under multiple runs.







### CONTACT INFORMATION

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